

CORP MODELLING USING FORMAL LANGUAGES

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**REZUMAT.** Modelarea corpurilor folosind limbajele formale. Definirea unui corp tridimensional se poate realiza de multe ori mai natural, și mai eficient utilizând comenzi de deplasare a "cursorului" (virf de creion imaginar) în diverse direcții. În acest fel se pot defini muchiile sau chiar suprafețele unui corp, prin "plimbarea" cursorului pe directii paralele cu axele de coordonate. Pentru aceasta am definit un limbaj de comandă și o serie de transformări elementare : translație, rotație, scalare, simetrie și proiecție pentru reprezentarea corpurilor descrise.

Solid corps modelling (tridimensional graphic objects) through front specification can be realised covering (describing) the edges ("wire-frame" method). This cover may be realised by moving commands of the cursor on a tridimensional frame this way: up, down, right, left, front, back. These moving commands compound the alphabet needed in the object outline description, which we want to model :

$$\Sigma = \{u, d, r, l, f, b\}$$

$w \in \Sigma^*$  is called command word of the cursor.

We will use for object definition an infinite frame of points obtained through intersection of the parallel plans with  $xOy$ ,  $xOz$  and  $yOz$  at distance equal to 1. Two points  $P(x,y,z)$  and  $P'(x',y',z')$  are "neighbours" if we have the following relation :

$$|x-x'| + |y-y'| + |z-z'| = 1 \quad (*)$$

This relation can be generalized :

$$x' = x + dx$$

$$y' = y + dy$$

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$$z' = z + dz$$

where  $dx, dy, dz \in \{-1, 0, 1\}$ ,  $|dx| + |dy| + |dz| > 0$ , which need an extension of the alphabet  $\Sigma$  by definition of other commands, obtained through composition of the existing commands.

We will use definition (\*) to simplify grammars.

One command word  $w \in \Sigma^*$  is, in fact, a succession of commands to the cursor and defines a "walk on the frame".

For a point  $P(x, y, z)$  we define its successors on the six directions this way :

- up :  $SUC((x, y, z), u) := (x, y, z+1)$
- down :  $SUC((x, y, z), d) := (x, y, z-1)$
- left :  $SUC((x, y, z), l) := (x, y-1, z)$
- right :  $SUC((x, y, z), r) := (x, y+1, z)$
- front :  $SUC((x, y, z), f) := (x+1, y, z)$
- back :  $SUC((x, y, z), b) := (x-1, y, z)$

The design specified in this way by a command word  $w \in \Sigma^*$  is defined as follows :

$$SUC : Z^3 \times \Sigma^* \dashrightarrow Z^3$$

$$SUC((x, y, z), \epsilon) = (x, y, z)$$

and

$$SUC((x, y, z), aw) = SUC(SUC(x, y, z), a), w)$$

We notice that the cursor movements are executed from left to right. The cursor movement may be defined initializing the cursor in  $P_0(x_0, y_0, z_0)$  and then moving through a command word  $w \in \Sigma^* : MOV(w) = SUC((x_0, y_0, z_0), w)$ .

A spatial graph unoriented described by

$$w = a_1 \dots a_n, (a_i \in \{u, d, l, r, f, b\}, i=1, n)$$

$$\text{is } g((x_0, y_0, z_0), w) = (V, A),$$

$$\text{where } V = \{ \text{MOV}(a_1 \dots a_i) / i=1, n \}$$

and

$$A = \{(\text{MOV}(a_1 \dots a_{i-1}), \text{MOV}(a_1 \dots a_i)) / i=1, n\}.$$

$$\text{We suppose that } \text{MOV}(\epsilon) = (x, y, z).$$

A language formed of command words is called language of solid object description, and the grammar which generates the language is called solid forms grammar.

To represent the objects described by such grammars we will define functions for spatial transformations.

#### TRANSLATION :

$$\begin{aligned} \text{TR}(g((x_0, y_0, z_0), w), (dx, dy, dz)) &= g((x_0 + dx, y_0 + dy, z_0 + dz), w) = \\ &= g((x_0, y_0, z_0), S_x^{|dx|} S_y^{|dy|} S_z^{|dz|} w) \end{aligned}$$

$$S_x = \begin{cases} f & dx > 0 \\ \epsilon & dx = 0 \\ b & dx < 0 \end{cases}$$

$$S_y = \begin{cases} l & dy < 0 \\ \epsilon & dy = 0 \\ r & dy > 0 \end{cases}$$

$$S_z = \begin{cases} u & dz > 0 \\ \epsilon & dz = 0 \\ d & dz < 0 \end{cases}$$

So, if  $V = (dx, dy, dz)$  is the translation vector then we may

say that  $\text{TR}(g(w), V) = g(S_x^{|dx|} S_y^{|dy|} S_z^{|dz|} w)$ .

For rotations with  $90^\circ$ ,  $180^\circ$ ,  $270^\circ$  or even with  $45^\circ$ ,  $135^\circ$ ,  $225^\circ$  and  $315^\circ$  we can apply command word transformations depending on the axis of the rotation Ox, Oy or Oz. For example, rotation around Ox will change the word w so :

$f$ ,  $b$  - remain unchanged

$u \rightarrow r$

$d \rightarrow l$

$l \rightarrow u$

$r \rightarrow d$ .

In the same way the substitutions applied to the productions rules of the grammar, for the other angles and axis are defined.

#### SCALATION :

A point  $P(x, y, z)$  through scalation become  $P'(x', y', z')$ , where:

$$\begin{aligned} |x'| &= f_x |x| \\ |y'| &= f_y |y| \\ |z'| &= f_z |z| \end{aligned}$$

If  $f_g f_x, f_g f_y, f_g f_z \geq 1$  then the image grows and we will consider these scalar factors natural.

If  $f_g f_x, f_g f_y, f_g f_z < 1$

$$f_g f_x = \frac{1}{f'_x}, \quad f_g f_y = \frac{1}{f'_y}, \quad f_g f_z = \frac{1}{f'_z} \text{ and} \\ f'_x, f'_y, f'_z \in \mathbb{N}$$

then the scalars applied to a command word  $w = a_1 \dots a_n$  will modify the word as follows :

$a_i \rightarrow a_i^f$ , where

$$f = \begin{cases} f_g \cdot f_z & \text{if } a_i \in \{u, d\} \\ f_g \cdot f_y & \text{if } a_i \in \{l, r\} \\ f_g \cdot f_x & \text{if } a_i \in \{f, b\} \end{cases}$$

These transformations may be applied to the grammars which generate the command language.

**SYMMETRIES :**

Reflections in planes  $xOy$ ,  $xOz$ ,  $yOz$  or in axis  $Ox$ ,  $Oy$ ,  $Oz$  are realised through such substitutions :

reflection in  $xOy$ :  $u \leftrightarrow d$

$xOz$ :  $l \leftrightarrow r$

$yOz$ :  $f \leftrightarrow b$ .

$SIM_{Ox}(g((x_0, y_0, z_0), w)) = g((-x_0, -y_0, z_0), w')$  , where  $w'$  is obtained by substituting  $l \leftrightarrow r$  and  $f \leftrightarrow b$ .

In the same way we construct reflections in  $Ox$  and  $Oy$ .

**PROJECTIONS :**

Projection parallel to  $Oz$  on  $xOy$  is defined as :

$$PR_{xOy}(g(w)) = \begin{cases} g(d^{z_0} w') & \text{if } z_0 > 0 \\ g(u^{z_0} w') & \text{if } z_0 < 0 \end{cases}$$

where  $w'$  is obtain by substituting " $u$ " or " $d$ " with  $\epsilon$ .

We can realise images with solid forms grammars extending the alphabet  $\Sigma$  with symbols representing commands to choose (select) colours, where the null colour , "dark", shift the cursor without effective translation :

$$\Sigma' = \Sigma \cup \{0, 1, \dots, 7\}.$$

In this case we consider that the initial position of the cursor is  $O(0,0,0)$  and the initial colour is null. Then the word

$$(S_x^{|x_0|} S_y^{|y_0|} S_z^{|z_0|} C w)$$

draw the object  $g(w)$  from the initial position  $(x_0, y_0, z_0)$  in the colour "c".

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