

## ALGEBRAIC SPECIFICATION OF PSP

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**REZUMAT.** - Specificarea algebrică a RSP. În lucrare se prezintă specificarea algebrică a tipurilor de date folosite în procesorul simbolic Poisson PSP ([Pâr89]). Se folosește un model ierarhic, bazat pe semantica denotațională (algebrică).

**0. Introduction.** Algebraic specification of PSP (Poisson Symbolic Processor, see [Pâr89]) data types is made by means of an hierachical model, from simple to complex types. The abstract data structure of PSP is viewed as a hierachy of abstract data types.

A Poisson series have the form:

$$S = \sum_{i=0}^{\infty} C_i x_1^{j_1} x_2^{j_2} \dots x_m^{j_m} \frac{\sin}{\cos} (k_1 y_1 + k_2 y_2 + \dots + k_n y_n), \quad (1)$$

where:  $C_i$  are numerical coefficients;  $x_1, x_2, \dots, x_m$  are monomial variables;  $y_1, y_2, \dots, y_n$  are trigonometric variables;  $j_1, j_2, \dots, j_m$  and  $k_1, k_2, \dots, k_n$  are exponents, and, respectively, coefficients. The summation index  $i$  covers the set of all possible combinations of the exponents  $j$  and coefficients  $k$  ( $j \in \mathbb{Z}^m, k \in \mathbb{Z}^n$ ).

The form (1) of Poisson series can be briefly written as follows:

$$S = \sum_{i=0}^{\infty} T_i, \quad (2)$$

in which  $T_i$  is a term of this series:

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$$T_i = C_i P_i F_i ,$$

where the polynomial part  $P_i$  has the form:

$$P_i = x_1^{j_1} x_2^{j_2} \dots x_m^{j_m}, \quad (3)$$

while the trigonometric part  $F_i$  is:

$$F_i = \frac{\sin}{\cos}(k_1 y_1 + k_2 y_2 + \dots + k_n y_n). \quad (4)$$

In practice, one does not operate with Poisson series, but with partial sums of these ones, called *Poisson expressions*, of the form:

$$S = \sum_{i=0}^N T_i, \quad N \in \mathbb{N}. \quad (5)$$

PSP operates with Poisson expressions of the form (5). The hierarchical model of its algebraic specification consists of five levels:

- 1) numerical coefficients specification;
- 2) trigonometric part specification;
- 3) polynomial part specification;
- 4) term specification;
- 5) Poisson expression specification.

**1. Numerical coefficients specification.** The coefficients  $C_i$  from (1) are considered rational numbers, of the form  $M/N$ , with  $M, N \in \mathbb{Z}$ . The definition of the abstract data type RAT follows the chain:

$$\text{BOOL} \text{ --> NAT --> INT --> RAT}$$

where:

BOOL - represents the primitive boolean type;

NAT - represents the hierarchical natural type (including zero value);

INT - represents the hierarchical integer type;

RAT - represents the hierarchical rational type.

In the specification of the NAT type we use the NAT\* subtype of NAT, which corresponds to the natural numbers without zero value. The above listed types are specified as follows:

**1.1. The primitive BOOL type**

```

type BOOL =
  SORT bool,
  CONS
    true: --> bool,
    false: --> bool,
  OPNS
    .not.: bool --> bool,
    .and.: bool, bool --> bool,
    .or.: bool, bool --> bool,
  VARS
    X, Y: bool,
  AXIOMS
    true * false,
    .not. true = false,
    .not. false = true,
    true .and. X = X,
    false .and. X = false,
    X .or. Y = .not.((.not. X) .and. (.not. Y))
endoftype.
  
```

**1.2. The NAT and NAT\* types**

The NAT type uses the primitive type BOOL:

```

type NAT =
  SORT bool, nat,
  CONS
    zero: --> nat,
    succ: nat --> nat,
    pred: (nat x: x noteq zero) --> nat,
  OPNS
    *: nat, nat --> nat,
    eq: nat, nat --> bool,
    +: nat, nat --> nat,
    noteq: nat, nat --> bool,
    < : nat, nat --> bool,
    [_,_]: nat, (nat x: x noteq zero) --> nat,
    GCD( _,_ ): nat, nat --> nat,
  VARS
    R, P, M, N: nat,
  AXIOMS
    * N eq M = M eq N,
  
```



```

*   zero eq zero = true,
*   zero eq succ(N) = false,
*   succ(N) eq succ(M) = N eq M,
*   N noteq M = not (N eq M),
*   pred(succ(N)) = N,
    [N,1] = N,
    ( $\exists R, P: \text{nat} : R < M, N = P * M + R$ ) ==> [N , M] = P,
    GCD(N, 0) = N,
    U < V ==> GCD(V, U) = GCD(U, V - U * [V , U]),
    GCD(U, V) = GCD(V, U),
*   N * 0 = 0 = 0 * N,
*   N * succ(M) = (N * M) + N = succ(M) * N,
*   N * pred(M) = (N * M) - N = pred(M) * N,
*   N + 0 = 0 + N = N,
*   N + succ(M) = succ(N + M) = succ(M) + N,
*   N + pred(M) = pred(N + M) = pred(M) + N,
    (0 < succ(0)) = true,
    (pred(N) < succ(N)) = true,
    (N < succ(N)) = true,
    (pred(N) < N) = true,
    N < M ==> (succ(N) < succ(M) = true),
    N < M ==> (pred(N) < pred(M) = true),
*   N - 0 = N,
*   IF M < N THEN N - succ(M) = pred(N - M),
*   IF M < N THEN N - pred(M) = succ(N - M),

```

endoftype.

We define NAT\*, subtype of the NAT type:

```

type NAT* =
  SORT bool, nat*,
  CONS
    one: --> nat*,
    succ: nat* --> nat*,
    pred: (nat* x: x noteq one) --> nat*,
  OPNS
    +: nat* --> nat*,
    -: nat* --> nat*,
    *: nat*, nat* --> nat*,
    eq: nat*, nat* --> bool,
    noteq: nat*, nat* --> bool,
  VARS
    M, N: nat*,
  AXIOMS
    The axioms of NAT type marked with * are axioms of the
    subtype NAT*, with the following modifications:
    one eq one = true,
    one eq succ(N) = false,
    N * 1 = 1 * N = N,
    N + 1 = 1 + N = succ(N),
    IF 1 < N THEN N - 1 = pred(N),

```

endoftype.

**1.3. The INT type**

The specification of INT type uses BOOL and NAT as primitive types:

```

type INT ≡
  SORT bool, int, nat,
  CONS
    zero: --> int,
    succ: int --> int,
    pred: int --> int,
  OPNS
    eq: int, int: --> bool,
    noteq: int, int --> bool,
    +: int, int --> int,
    -: int, int --> int,
    *: int, int --> int,
    <: int, int --> bool,
    |_ |: int --> nat,
  VARS
    N, M: int,
  AXIOMS
    N eq M = M eq N,
    zero eq zero = true,
    succ(N) eq succ(M) = N eq M,
    N noteq M = not (N eq M),
    pred(succ(N)) = N,
    succ(pred(N)) = N,
    | 0 | = 0,
    (pred(0) < 0) = true,
    The axioms of INT which refer to the operations +, -
    and < are inherited from NAT type;
endofstype.

```

*Remark.* The INT type contains the zero value and two unary operations, *succ* and *pred*. The integer number  $n$  ( $n > 0$ ) is obtained by  $n$  successive applications of the *succ* operation to 0. Analogously, the integer number  $-n$  ( $n > 0$ ) is constructed with *pred*, starting with 0.

**1.4. The RAT type**

The specification of RAT type uses INT and NAT\* as primitive types:

```

type RAT ≡
  SORT int, nat*, rat,
  CONS

```

```

    _/_: int, nat* --> rat,
OPNS  +_: rat, rat --> rat,
      *_: rat, rat --> rat,
      -: rat, rat --> rat,
      _ -: rat, rat --> rat,
VARS  M/N, P/Q: rat,
      S, T: nat*,
AXIOMS
      (M/N) + (P/Q) = ((M * Q) + (P * N)) / (N * Q),
      (M/N) * (P/Q) = (M * P) / (N * Q),
      (M/N) : (P/Q) = (M * Q) / (N * P),
      (M/N) - (P/Q) = ((M * Q) - (P * N)) / (N * Q),
      (∃ S: nat*, ∃ T: nat*, :
        M = S * GCD(|M|, |N|) ∧ N = T * GCD(|M|, |N|)) -
        → M/N = S/T

```

endoftype.

**2. Trigonometric part specification.** The definition of the trigonometric parts  $F_1$  from (4) follows the chain:

SET --> FLIN --> TTR

where:

SET - represents a symbol set;

FLIN - represents the abstract type of linear forms;

TTR - represents the abstract type of trigonometric parts.

**2.1. The SET type**

The SET type consists of a set of symbols. This type is used as primitive type for the specification of FLIN type.

```

type SET =
  SORT set,
  CONS
    X1 : --> set,
    X2 : --> set,
    - - - - -
    Xn : --> set,
    Y1 : --> set,
    Y2 : --> set,
    - - - - -
    Ym : --> set

```

endoftype.

### 2.2. The FLIN type

The FLIN type defines the linear forms over the types SET and INT.

```

type FLIN =
  SORT int, set, flin,
  CONS
    _·_ (concatenation) : int, set --> flin,
  OPNS
    _+_ : flin, flin --> flin,
    _-_- : flin, flin --> flin,
    _*_ : int, flin --> flin,
  VARS
    X, Y, Z: set,
    M, N, K: int,
  AXIOMS
    X·1 = 1·X = X,
    X·0 = 0·X = 0,
    N·X = X·N,
    X·(N + M) = (N + M)·X = N·X + M·X,
    N·X + M·Y = M·Y + N·X,
    (N·X + M·Y) + K·Z = N·X + (M·Y + K·Z),
    N·X + 0 = 0 + N·X = N·X,
    N * 0 = 0 = 0 * N,
    1 * X = X * 1 = X,
    N * (M·X) = (N * M)·X,
    N * (M·X ± K·Y) = (N * M)·X ± (M * K)·Y
endofype.

```

The FLIN type contains linear combinations of the form:

$$N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k$$

which can be used as arguments for sine and cosine functions in trigonometric parts  $F_i$ . The three operations (+, \* and concatenation), together with the axioms, are used for the specification of the trigonometric part.

### 2.3. The TTR type

```

type TTR =
  SORT flin, ttr,
  CONS
    cos: flin --> ttr,
    sin: flin --> ttr,
  AXIOMS
    cos 0 = 1,
    sin 0 = 0,
endofype.

```

**3. Polynomial part specification.** The definition of the polynomial parts  $P_i$  from (3) follows the scheme:

MPOL --> PPOL

where:

MPOL - represents the set of monomials;

PPOL - represents the set of polynomials.

**3.1. The MPOL type**

The MPOL type defines the set of monomials of the form  $X^N$ , where  $X$  is in SET and  $N$  is of INT type.

```

type MPOL =
  SORT int, set, mpol,
  CONS
    _ ^ _ (raise to power) : set, int --> mpol
              ( X^N = X^N )
  VARS
    X: set, N: int,
  AXIOMS
    (X ^ 0) = 1,
    (X ^ 1) = X,
endoftype.

```

**3.2. The PPOL type**

The PPOL type defines the set of polynomials of the form (3):

```

type PPOL =
  SORT mpol, ppol
  CONS
    _ * _ : mpol, mpol --> mpol,
  VARS
    X^M , X^N , Y^M , Z^K : ppol,
  AXIOMS
    X^M * X^N = X^N * X^M = X^{M+N} ,
    X^N * Y^M = Y^M * X^N ,
    X^N * 1 = X^N = 1 * X^N ,
    X^N * (Y^M * Z^K) = (X^N * Y^M) * Z^K
endoftype.

```



**4. Term specification.** The terms  $T_i$  from (2) are defined as follows:

```

type TERM ≡
  SORT rat, ttr, ppol, term,
  CONS
    · (concatenation) : rat, ttr, ppol --> term,
  VARS
    N/M : rat, cosY, sinY: ttr, X: ppol,
  * * * We denote with Y the linear form and with X the polynomial
  * * * form
  AXIOMS
      sinY                sinY
    (N/M · {   }) · X = N / M · ({   }) · X ,
      cosY                cosY
      sinY                sinY                sinY
    N/M · {   } · X = {   } · N/M · X = {   } · X · N/M
      cosY                cosY                cosY
                        = X · N/M · {   } ,
                        cosY
    0 · 0 · X = 0 ,
    1/1 · 1 · X = X ,
      sinY
    0 · {   } · X = N/M · 0 · X = 0
      cosY
endofdtype.

```

**5. Poisson expression specification.** Taking into account the above definitions, the Poisson expressions (5) will be defined in the following form:

```

type EXP ≡
  SORT term, exp,
  CONS
    +_ : term, term --> exp,
    -_ : term, term --> exp,
    *_ : term, term --> exp,
  OPNS
    ∂ (differentiation) : exp, set --> exp,
    ∫ (integration) : exp, set --> exp,
  VARS
    N/M , P/Q : rat,
    Y, Y1 , Y2 : flin,
    X, X1 , X2 : mpol,
  AXIOMS
      sinY                sinY
    N/M · {   } · X ± P/Q · {   } · X =
      cosY                cosY

```

$$\begin{aligned}
 &= (N/M \pm P/Q) \cdot X \cdot \left\{ \frac{\sin Y}{\cos Y} \right\}, \\
 (N/M \cdot \cos Y_1 \cdot X_1) * (P/Q \cdot \sin Y_2 \cdot X_2) &= \\
 &= (1/2 * N/M * P/Q) \cdot X_1 \cdot X_2 \cdot \sin(Y_1+Y_2) + \\
 &+ (1/2 * N/M * P/Q) \cdot X_1 \cdot X_2 \cdot \sin(Y_2-Y_1), \\
 (N/M \cdot \sin Y_1 \cdot X_1) * (P/Q \cdot \sin Y_2 \cdot X_2) &= \\
 &= (1/2 * N/M * P/Q) \cdot X_1 \cdot X_2 \cdot \cos(Y_1-Y_2) - \\
 &- (1/2 * N/M * P/Q) \cdot X_1 \cdot X_2 \cdot \cos(Y_1+Y_2), \\
 (N/M \cdot \cos Y_1 \cdot X_1) * (P/Q \cdot \cos Y_2 \cdot X_2) &= \\
 &= (1/2 * N/M * P/Q) \cdot X_1 \cdot X_2 \cdot \cos(Y_1+Y_2) + \\
 &+ (1/2 * N/M * P/Q) \cdot X_1 \cdot X_2 \cdot \cos(Y_1-Y_2),
 \end{aligned}$$

$$\frac{\partial}{\partial Y_k} (N/M \cdot \left\{ \frac{\cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1)}{\sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1)} \right\} \cdot X_1^{M_1} \cdot \dots \cdot Y_k^{M_k} \cdot \dots \cdot X_h^{M_h}) =$$

$$= N \cdot M_k / M \cdot \left\{ \frac{\cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1)}{\sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1)} \right\} \cdot X_1^{M_1} \cdot \dots \cdot Y_k^{M_k-1} \cdot \dots \cdot X_h^{M_h} \mp$$

$$\mp N \cdot N_k / M \cdot \left\{ \frac{\sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1)}{\cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1)} \right\} \cdot X_1^{M_1} \cdot \dots \cdot Y_k^{M_k} \cdot \dots \cdot X_h^{M_h},$$

$$I_0^{N_k} = \int N/M \cdot \sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot$$

$$\cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^0 \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} dY_k =$$

$$= (-1/N_k \cdot N/M) \cdot \cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot$$

$$\cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^0 \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot Y_h^{M_h}$$

$$I_1^{N_k} = \int N/M \cdot \sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot$$

$$\cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^1 \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} dY_k =$$

$$\begin{aligned}
 &= (-1/N_k * N/M) \cdot \cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^1 \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} + \\
 &+ (1/(N_k * N_k)) \cdot \sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^0 \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} \\
 I_p^{N_k} &= \int N/M \cdot \sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^p \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} dY_k = \\
 &= (-1/N_k * N/M) \cdot \cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^{M_k} \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} + \\
 &+ (N/M * p / (N_k * N_k)) \cdot \sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^{p-1} \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} - \\
 &\quad - (N/M * p / N_k * (p-1) / N_k) \cdot I_{p-2}^{N_k} \\
 \text{endoftype.}
 \end{aligned}$$

Remark. The EXP type must also contain the axioms referring to the recurrent computation of the integral:

$$\begin{aligned}
 J_p^{M_k} &= \int N/M \cdot \cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^p \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} dY_k = \\
 &= (-1/N_k * N/M) \cdot \sin(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^p \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} + \\
 &+ (N/M * p / (N_k * N_k)) \cdot \cos(N_1 \cdot Y_1 + N_2 \cdot Y_2 + \dots + N_k \cdot Y_k + \dots + N_1 \cdot Y_1) \cdot \\
 &\quad \cdot X_1^{M_1} \cdot \dots \cdot X_{k-1}^{M_{k-1}} \cdot Y_k^{p-1} \cdot X_{k+1}^{M_{k+1}} \cdot \dots \cdot X_h^{M_h} - \\
 &\quad - (N/M * p / N_k * (p-1) / N_k) \cdot J_{p-2}
 \end{aligned}$$

CONCLUDING REMARKS. This paper specifies the data types

defined in PSP (implemented in Pascal) in a hierarchical way. Some of the advantages of algebraic specifications are used, especially those involving the correctness of the defined operations. In such a way the singular cases (as non-determinations or exceptions) are avoided. The authors intention is to simulate the operational semantics of PSP by using terms rewrite rules. PSP can also be specified in the specification and programming language OBJ3 (see [Gog88] and [Kir87]).

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