

Admission Exam September 2018
Written Exam for Computer Science

Important observations for candidates:

1. All arrays are indexed starting from 1.
2. Multiple answer problem statements (Part A) can have one or more correct answers. Candidates must write the answers on their contest sheet (not the sheets with the problem statements). Solutions are graded only if they identify all the correct answers, and only them.
3. Complete solutions are expected for problem statements in Part B.
 - a. Solutions will be described using *pseudocode* or a *programming language* (Pascal/C/C++).
 - b. The first criterion when evaluating solutions is algorithm *correctness*, followed by its *performance* regarding *time of execution* and *required memory space*.
 - c. *Providing description and justification* for all (sub)algorithms before their implementation is *mandatory*. In addition, *comments* will be used to describe details of the technical implementation, meaning of used identifiers, data structures and so on. Failure to observe these requirements leads to a 10% penalty from the problem score.
 - d. Using predefined functions or libraries is prohibited (for example: *STL*, predefined string functions).

Part A (30 points)

A.1. What does the following subalgorithm do? (5p)

The generate(n) subalgorithm processes natural number n ($0 < n < 100$).

```
Subalgorithm generate(n):
  nr ← 0
  For i ← 1, 1801 do
    usedi ← false
  EndFor
  While not usedn do
    sum ← 0, usedn ← true
    While (n ≠ 0) do
      digit ← n MOD 10, n ← n DIV 10
      sum ← sum + digit * digit * digit
    EndWhile
    n ← sum, nr ← nr + 1
  EndWhile
  return nr
EndSubalgorithm
```

State the effect of this subalgorithm.

- A. repeatedly calculates the sum of the cubes of n 's digits until the sum is equal to n , after which it returns the number of repetitions
- B. calculates the sum of the cubes of n 's digits and returns this sum
- C. calculates the sum of the cubes of n 's digits, replaces n with the obtained sum and returns it
- D. repeatedly calculates the sum of the cubes of n 's digits until a sum is obtained for the second time and returns the number of repetitions
- E. calculates how many times number n was replaced with the sum of the cubes of n 's digits until a previously calculated value or the number itself is obtained; this number is then returned

A.2. What values are required? (5p)

Consider subalgorithm process(v , k), where v is an array of k natural numbers ($1 \leq k \leq 1\ 000$).

```
Subalgorithm process(v, k)
  i ← 1, n ← 0
  While i ≤ k and vi ≠ 0 do
    y ← vi, c ← 0
    While y > 0 do
      If y MOD 10 > c then
        c ← y MOD 10
      EndIf
      y ← y DIV 10
    EndWhile
    n ← n * 10 + c
    i ← i + 1
  EndWhile
  return n
EndSubalgorithm
```

State for which values of v and k the subalgorithm returns 928.

- A. $v = (194, 121, 782, 0)$ and $k = 4$
- B. $v = (928)$ and $k = 1$
- C. $v = (9, 2, 8, 0)$ and $k = 4$
- D. $v = (8, 2, 9)$ and $k = 3$
- E. $v = (912, 0, 120, 8, 0)$ and $k = 5$

A.3. Logical Evaluation (5p)

Let us consider array s with k boolean elements and subalgorithm $\text{evaluation}(s, k, i)$, where k and i are natural numbers ($0 \leq i \leq k \leq 100$).

```

Subalgorithm evaluation(s, k, i)
  If i ≤ k then
    If si then
      return si
    else
      return (si or evaluation(s, k, i + 1))
    EndIf
  else
    return false
  EndIf
EndSubalgorithm

```

State how many times subalgorithm $\text{evaluation}(s, k, i)$ is called itself given the following instruction sequence:

```

s ← (false, false, false, false, false, false, true, false, false, false)
k ← 10
i ← 3
evaluation(s, k, i)

```

- A. 3 times
- B. the same number of times as in the following instruction sequence

```

s ← (false, false, false, false, false, false, false, true)
k ← 8
i ← 4
evaluation(s, k, i)

```

- C. 6 times
- D. never
- E. infinite number of times

A.4. Reunion (5p)

Consider subalgorithm $\text{belongs}(x, a, n)$, which checks whether natural number x belongs to set a having n elements; a is an array of n elements that represents a natural numbers set ($1 \leq n \leq 200$, $1 \leq x \leq 1000$). Let us consider subalgorithms $\text{reunion}(a, n, b, m, c, p)$ and $\text{compute}(a, n, b, m, c, p)$, described below, where a, b and c are arrays that represent natural number sets having n, m and p elements respectively ($1 \leq n \leq 200$, $1 \leq m \leq 200$, $1 \leq p \leq 400$). Input parameters are a, n, b, m and p , and output parameters are c and p .

<pre> 1. Subalgorithm reunion(a, n, b, m, c, p): 2. If n = 0 then 3. For i ← 1, m do 4. p ← p + 1 5. c_p ← b_i 6. EndFor 7. else 8. If not belongs(a_n, b, m) then 9. p ← p + 1 10. c_p ← a_n 11. EndIf 12. reunion(a, n - 1, b, m, c, p) 13. EndIf 14. EndSubalgorithm </pre>	<pre> 1. Subalgorithm compute(a, n, b, m, c, p): 2. p ← 0 3. reunion(a, n, b, m, c, p) 4. EndSubalgorithm </pre>
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Which of the following statements are always true:

- A. when set a contains a single element, calling subalgorithm $\text{compute}(a, n, b, m, c, p)$ results in an infinite loop
- B. when set a contains 4 elements, calling subalgorithm $\text{compute}(a, n, b, m, c, p)$ results in executing the instruction on line 12 a number of 4 times
- C. when set a contains 5 elements, calling subalgorithm $\text{compute}(a, n, b, m, c, p)$ results in executing the instruction on line 2 of the reunion subalgorithm a number of 5 times
- D. when sets a and b have the same number of elements, after the execution of subalgorithm $\text{compute}(a, n, b, m, c, p)$ set c will have the same number of elements as set a
- E. when sets a and b have the same elements, after the execution of subalgorithm $\text{compute}(a, n, b, m, c, p)$ set c will have the same number of elements as set a

A.5. Exponentiation (5p)

Which of the following algorithms correctly computes a^b , where a and b are natural numbers ($1 \leq a \leq 11, 0 \leq b \leq 11$).

<p>A. Subalgorithm expo(a, b): result ← 1 While b > 0 do If b MOD 2 = 1 then result ← result * a EndIf b ← b DIV 2 a ← a * a EndWhile return result EndSubalgorithm</p>	<p>B. Subalgorithm expo(a, b): If b ≠ 0 then If b MOD 2 = 1 then return expo(a * a, b / 2) * a else return expo(a * a, b / 2) EndIf else return 1 EndIf EndSubalgorithm</p>
<p>C. Subalgorithm expo(a, b): result ← 1 While b > 0 do result ← result * a b ← b - 1 EndWhile return result EndSubalgorithm</p>	<p>D. Subalgorithm expo(a, b): If b = 0 then return 1 EndIf aux ← expo(a, b DIV 2) If b MOD 2 = 0 then return aux * aux else return a * aux * aux EndIf EndSubalgorithm</p>
<p>E. Subalgorithm expo(a, b): If b = 0 then return 1 EndIf return a * expo(a, b - 1) EndSubalgorithm</p>	

A.6. Largest multiple (5p)

Which of the following subalgorithms returns the largest multiple of number a , multiple that is smaller or equal with natural number b ($0 < a < 10\,000, 0 < b < 10\,000, a < b$)?

<p>A. Subalgorithm f(a, b): c ← b While c MOD a = 0 do c ← c - 1 EndWhile return c EndSubalgorithm</p>	<p>B. Subalgorithm f(a, b): If a < b then return f(2 * a, b) else If a = b then return a else return b EndIf EndIf EndSubalgorithm</p>
<p>C. Subalgorithm f(a, b): return (b DIV a) * a EndSubalgorithm</p>	<p>E. Subalgorithm f(a, b): c ← a While c < b do c ← c + a EndWhile If c = b then return c else return c - a EndIf EndSubalgorithm</p>
<p>D. Subalgorithm f(a, b): If b MOD a = 0 then return b EndIf return f(a, b - 1) EndSubalgorithm</p>	

Part B (60 points)

B.1. Polynomial evaluation (10 points)

Consider the evaluation(n , coef, x) subalgorithm, where *coef* is a vector containing $n + 1$ real numbers in interval $[-100, 100]$ which represent the coefficients of the n degree polynome $P(x) = \text{coef}_1 * x^n + \text{coef}_2 * x^{n-1} + \dots + \text{coef}_n * x + \text{coef}_{n+1}$, provided in decreasing order of the powers of x (n is a natural number, $1 \leq n \leq 10$). The subalgorithm calculates the polynome's value at point x (x is a real number in interval $[-10, 10]$).

```
Subalgorithm evaluation(n, coef, x):
  val ← 0.0
  For i ← 1, n + 1 do
    val ← val * x + coef[i]
  EndFor
  return val
EndSubalgorithm
```

Write a *recursive* implementation without loops for subalgorithm evaluation(n , coef, x), which has the same signature and effect.

B.2. Intersection (25 points)

Let us consider two arrays, each containing *distinct* integer numbers between -30 000 and 30 000. Array a has n ($0 < n \leq 10\,000$) elements, and array b has m ($0 < m \leq 10\,000$) elements *sorted in ascending order*.

Write a subalgorithm that determines array c having k ($0 \leq k \leq 10\,000$) elements formed using all the *common* elements of the two arrays, each of them appearing a single time, in any order. The input parameters are the two arrays (a and b) and their lengths (n and m). The array c and its length k are the output parameters. If there are no common elements, then k is 0.

Example: if $n = 4$, $a = (5, -7, -2, 3)$, $m = 5$ and $b = (-2, 3, 5, 7, 8)$, array c has $k = 3$ elements and $c = (5, -2, 3)$.

B.3. Brotherless number sequence (25 points)

Let us consider array x with n ($0 < n \leq 10\,000$) *distinct*, non-zero natural numbers smaller than 30 000. Two numbers are called *brothers* if they are *distinct* and if they have *at least two distinct digits in common*. For example, 5867 and 17526 are *brothers*, but 5867 and 152 are not *brothers*. Also, 131 and 114 are not *brothers*.

Requirements:

- i. Write a subalgorithm that checks whether natural numbers a and b are brothers ($0 < a \leq 30\,000$, $0 < b \leq 30\,000$). Input parameters are the two numbers a and b . The output parameter *isBrother* will be *true* if a and b are brothers, and *false* otherwise.
- ii. Write a subalgorithm that determines the longest subarray of x that consists of elements that *have no brother* in array x . A subarray of an array is made up of array's elements in consecutive positions. Input parameters are array x and its length n . Output parameters are the starting index of the subarray *start* and its length k . In case multiple subarrays of the same length exist, consider the last one of them. If no such subarray exists, then *start* will be -1 and k will be 0.

Example: If $n = 11$ and $x = (12345, 9, 100, 567, 5678, 345, 123, 8989, 222, 11, 78)$. The brotherless numbers contained in x are 9, 100, 8989, 222, 11, and the longest subarray is (8989, 222, 11), so *start* = 8 and $k = 3$.

Note:

1. All subjects are mandatory.
2. Drafts are not taken into consideration for grading.
3. Default 10 points.
4. Working time is 3 hours.

BAREM

OFICIU 10 puncte

Partea A 30 puncte

- A. 1. Oare ce face? Răspunsul E..... 5 puncte
- A. 2. Ce valori sunt necesare? Răspunsurile A, C 5 puncte
- A. 3. Evaluare logică. Răspunsul B 5 puncte
- A. 4 Reuniune. Răspunsurile B, E 5 puncte
- A. 5. Exponențiere. Răspunsurile A, B, C, D, E 5 puncte
- A. 6. Cel mai mare multiplu. Răspunsurile C, D, E..... 5 puncte

Partea B 60 puncte

B. 1. Evaluare polinom 10 puncte

- respectarea parametrilor de intrare și ieșire..... 2 puncte
- condiția de oprire din recursivitate 2 puncte
- autoapel 2 puncte
- valoarea returnată la oprirea recursivității 2 puncte
- valoarea returnată la continuarea recursivității 2 puncte

B. 2. Intersecție 25 puncte

- respectarea parametrilor de intrare și ieșire..... 2 puncte
 - variante:
 - folosirea căutării binare 23 puncte
 - fără căutare binară maxim 18 puncte

B. 3. Secvență de numere fără frați 25 puncte

- respectarea parametrilor de intrare și ieșire..... 2 puncte
- proprietatea de frate..... 10 puncte
- determinarea unei secvențe 9 puncte
- determinarea celei mai lungi secvențe..... 4 puncte